

# Lecture #34: Review of Scheme

# List Tails

- The full Scheme library includes the function `list-tail`, which essentially performs the `cdr` operation a given number of times:

```
scm> (list-tail '(a b c d) 0)
```

```
(a b c d)
```

```
scm> (list-tail '(a b c d) 1)
```

```
(b c d)
```

```
scm> (list-tail '(a b c d) 2)
```

```
(c d)
```

```
scm> (list-tail '(a b c d) 4)
```

```
()
```

- Can you implement it?

```
(define (list-tail lst k)
```

```
)
```

- Solutions to problems in this lecture are in `34.scm`.

# List Elements by Index

- Another Standard Scheme library operation is `list-ref`, which works like `lst[k]` in Python:

```
scm> (list-ref '(a b c d) 0)
```

```
a
```

```
scm> (list-ref '(a b c d) 3)
```

```
d
```

- What's the simplest implementation you can come up with?

```
(define (list-ref lst k)
```

```
)
```

# A Faster List?

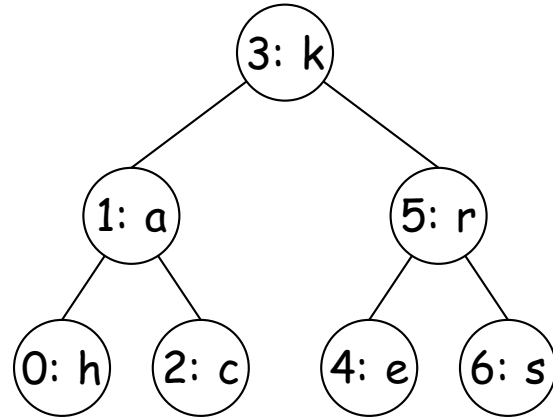
- Linked lists are problematic if your algorithm calls for performing lots of `list-ref` operations.
- In Python, indexed list access (like `A[i]`) takes constant ( $\Theta(1)$ ) time.
- What about `list-ref`?

# A Faster List?

- Linked lists are problematic if your algorithm calls for performing lots of `list-ref` operations.
- In Python, indexed list access (like `A[i]`) takes constant ( $\Theta(1)$ ) time.
- What about `list-ref`?  
It takes  $\Theta(N)$  time (worst case) for a list of length  $N$ .
- While we can't get to  $\Theta(1)$  with a linked list, we can do better than  $\Theta(N)$ .

# Lists as Trees

- One idea is to represent a list as a *binary search tree*.
- The labels of this tree contain integer indices (0- $N$ ) and the values at those indices.
- For example, (h a c k e r s) could be



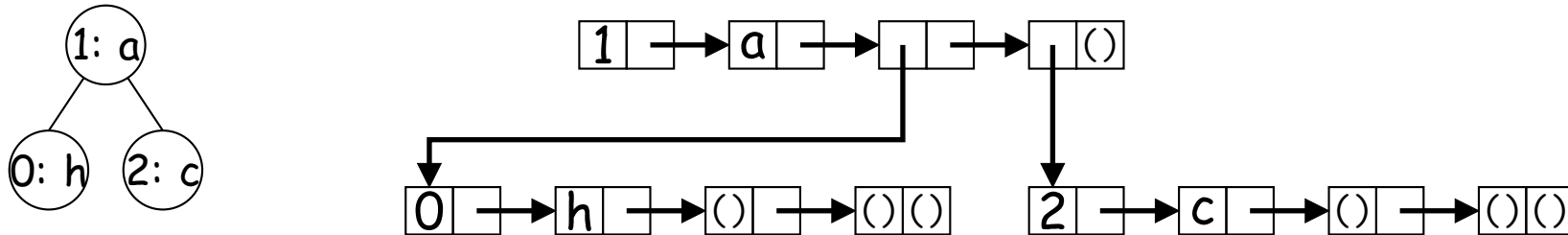
- So for any node in this tree, all the nodes in its left subtree have smaller indices, and all the nodes in the right subtree have larger indices.

# Lists as Trees as Lists

- First, let's define a suitable data structure to represent these trees.  
Suggestions?

# Lists as Trees as Lists

- First, let's define a suitable data structure to represent these trees.
- Each node can be represented as a list containing the index, value, and two children, with () representing an empty tree (as well as an empty list).

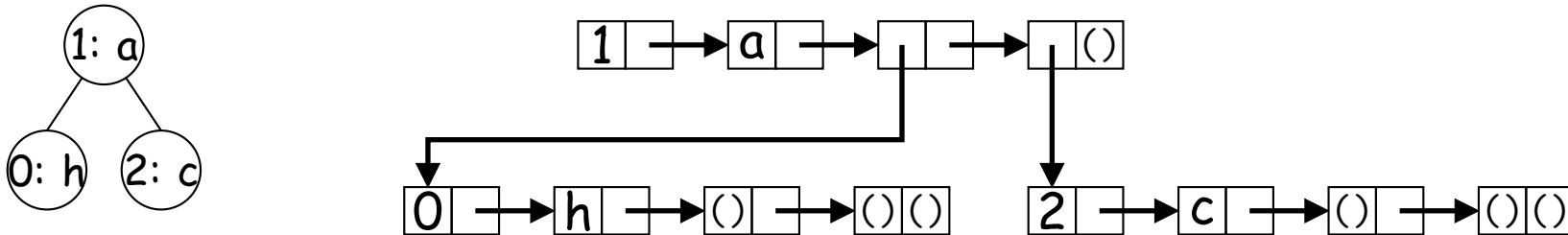


- What's a reasonable set of functions for accessing or creating these "array trees"?



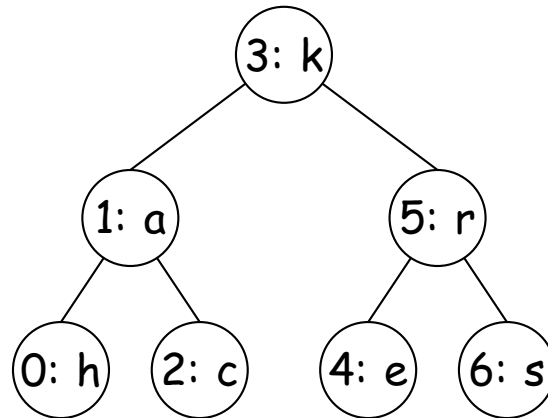
# Lists as Trees as Lists

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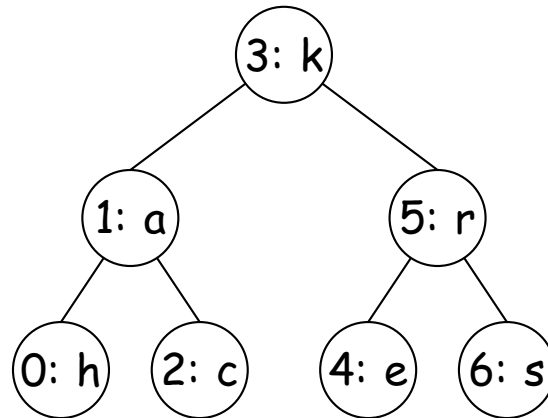
```
(define (arr-index arr) (car arr))
(define (arr-value arr) (car (cdr arr)))
(define (arr-left arr) (car (cdr (cdr arr))))
(define (arr-right arr) (car (cdr (cdr (cdr arr)))))
(define (arr-make index value left right)
  (list index value left right))
(define arr-empty nil)
```

# Fetching from an Array Tree



- To fetch item  $\#j$ , we start at the root and compare  $j$  to the index we find there. If equal, we return the value in that node.
- If  $j$  is less than the node's index, we search for item  $j$  in the left subtree. Otherwise, we search the right.
- How long does this process take if there are  $N$  elements in the represented array?

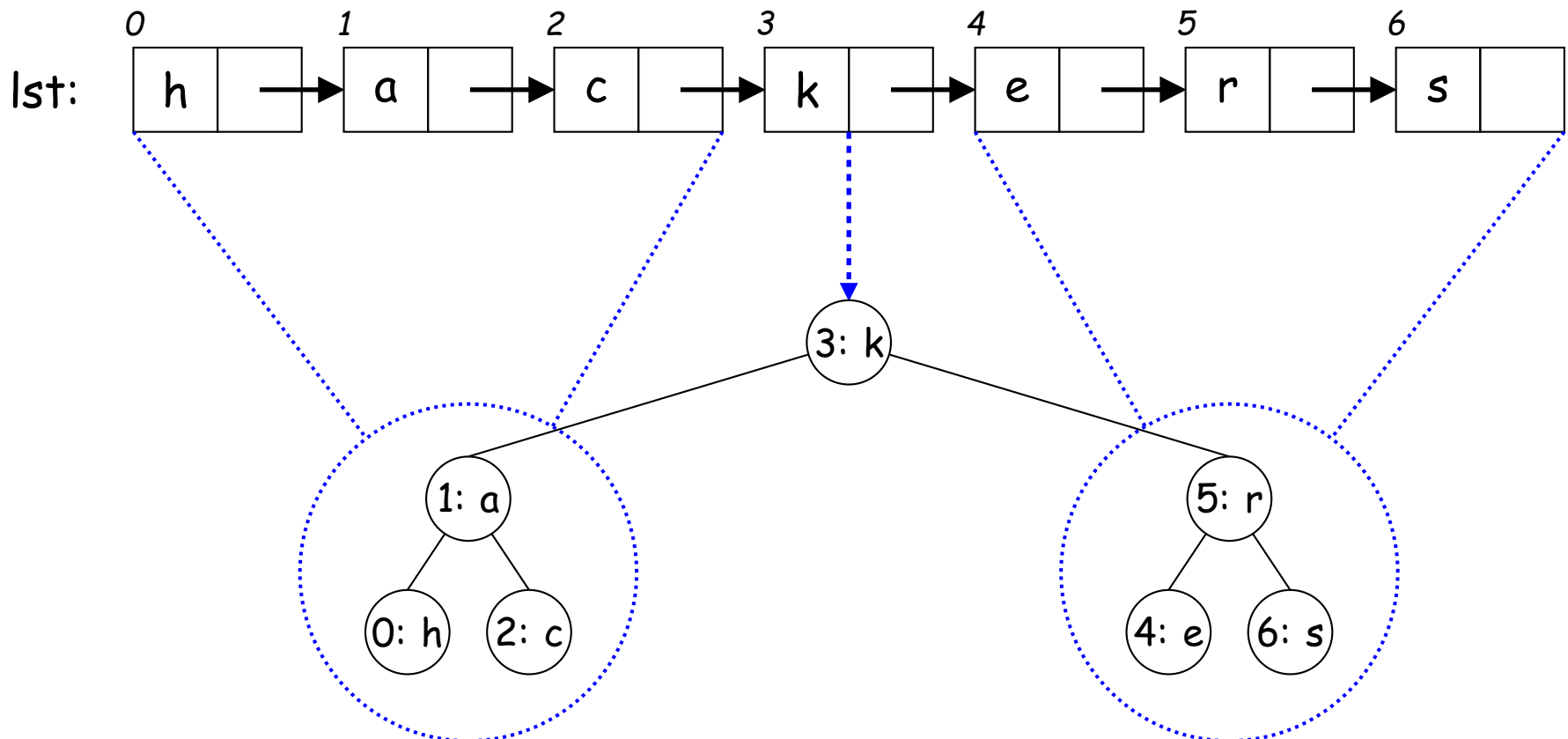
# Fetching from an Array Tree



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- If  $j$  is less than the node's index, we search for item  $j$  in the left subtree. Otherwise, we search the right.
- How long does this process take if there are  $N$  elements in the represented array?  
 $\Theta(\lg N)$  (assuming that the tree is evenly balanced.)
- How would we code it in Scheme?

# Building an Array Tree

- Given an arbitrary Scheme list, what is a function that will return an array tree such as we've been discussing?
- Suggested approach: find list's length, then build tree from middle element of list, with subtrees found recursively on either side.  
How do we do this in Scheme?



# Further Encapsulation

- Now that we have a function for creating this data structure, how do we use it to get the following effects?

```
scm> (define arr (make-array-tree '(h a c k e r s)))
```

```
arr
```

```
scm> (arr 0)
```

```
h
```

```
scm> (arr 4)
```

```
e
```

- In other words, create a function with a single (index) argument that returns the indexed item of the original list.

# Let\*

- Let's consider a macro.
- You implemented the `let` form, which should have the following behavior:

```
scm> (define x 4)
x
scm> (let ((x 5) (y x)) (list x y))
(5 4)
```

- That is, `x` in `(y x)` still refers to the outer `x`, not to its redefinition.
- There is another form in standard Scheme: `let*`, which is incremental:

```
scm> (let* ((x 5) (y x)) (list x y))
(5 5)
```

- ...as if we had written

```
scm> (let ((x 5))
...>   (let ((y x))
...>     (list x y)))
(5 5)
```

- Assuming that `let*` bodies have a single expression, how can we define `let*` to get this effect?