1 Scheme

1.1 What would Scheme do?

scm> (and 0 2 200)

200

scm> (or True (/ 1 0))

True

scm> (and False (/ 1 0))

False

scm> (not 3)

False

1.2 What would Scheme display?

scm> (define a (+ 1 2))

a

scm> a

3

scm> (define b (+ (+ 3 3) (* 4 4)))

b

scm> (+ a b)

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2  Scheme

scm> (= (modulo 10 3) (quotient 5 3))

#t

scm> (even? (+ (- (* 5 4) 3) 2))

#f

scm> (if (and #t (/ 1 0)) 1 (/ 1 0))

Error

scm> (if (> (+ 2 3) 5) (+ 1 2 3 4) (+ 3 4 (* 3 2)))

13

scm> ((if (< 9 3) + -) 4 100)

-96

scm> (if 0 #t #f)

#t

1.3 Write two Scheme expressions that are equivalent to the following Python statement - one defining a function directly, and the other creating an anonymous lambda that is then bound to the name cat:

```
cat = lambda meow, purr: meow + purr
```

```
(define cat (lambda (meow purr) (+ meow purr)))
(define (cat meow purr) (+ meow purr))
```

1.4 Spot the bug(s). Test out the code and your fixes in the scheme interpreter!

(https://scheme.cs61a.org/)

```
(define (sum-every-other lst)
  (cond ((null? lst) lst)
        (else (+ (cadr lst)
                 (sum-every-other (caddr lst))))))
```

1. Missing a paren at the end.
2. The base case should return 0, not `('.
3. (cadr lst) is a list, so it doesn’t make sense to add it to something. Instead, use
(car lst), which will give us a number.
4. Using the caar (car of the car) is incorrect because the car is a number and it
doesn’t make sense to get the car of a number. Instead, we should use cddr (the
cdr of the cdr) to skip forward two elements. However, the cdr could be '(), so we
need to add a case to our cond to take care of this.
The corrected function:

(define (sum-every-other lst)
  (cond ((null? lst) 0)
        ((null? (cdr lst)) (car lst))
        (else (+ (car lst)
                  (sum-every-other (cddr lst))))))

1.5 Define sixty-ones, a function that takes in a list and returns the number of times
that 1 follows 6 in the list.
>
(sixty-ones '(4 6 1 6 0 1))
1
(sixty-ones '(1 6 1 4 6 1 6 0 1))
2
(sixty-ones '(6 1 6 1 4 6 1 6 0 1))
3

(define (sixty-ones lst)
  (cond ((or (null? lst) (null? (cdr lst))) 0)
        ((and (= 6 (car lst)) (= 1 (cadr lst)))
         (+ 1 (sixty-ones (cddr lst))))
        (else (sixty-ones (cdr lst)))))

1.6 Define no-elevens, a function that takes in a number n, and returns a list of all
distinct length-n lists of 1s and 6s that do not contain two consecutive 1s.
>
(no-elevens 2)
((6 6) (6 1) (1 6))
(no-elevens 3)
((6 6 6) (6 6 1) (6 1 6) (1 6 6) (1 6 1))
(no-elevens 4)
((6 6 6 6) (6 6 6 1) (6 6 1 6) (6 1 6 6) (6 1 6 1) (1 6 6 6) (1 6 6 1) (1 6 1 6))

(define (no-elevens n)
  (cond ((= 0 n) '(()))
        ((= 1 n) '((6) (1)))
        (else (append (add-to-all 6 (no-elevens (- n 1)))
                      (add-to-all 1
                                 (add-to-all 6 (no-elevens (- n 2))))))))
1.7 Define `remember`, a function that takes in another zero-argument function `f`, and returns another function `g`. When called for the first time, `g` will call `f` and pass on its return value. When called subsequent times, `g` will remember its previous return value and return it directly, without calling `f` again.

(Hint: look up `set!` in the Scheme spec!)

```
(define (remember f)
  (define remembered? #f)
  (define remembered nil)
  (lambda ()
    (if remembered?
        remembered
        (begin (set! remembered (f))
               (set! remembered? #t)
               remembered)))
)
```

```
scm> (define (f) (print "hello!") 5)
scm> (define g (remember f))
scm> (f)
hello!
5
scm> (g)
hello!
5
scm> (g)
5
```

Check your understanding

- How are call expressions (like `(+ 1 2 3)`) evaluated? What about special forms, like `(or #f #t (/ 1 0))`?

  To evaluate call expressions, Scheme first evaluates the operator, and then evaluates all of the operands from left to right. It then applies the operator to the operands (i.e. calls the procedure with the evaluate operands), just like how Python evaluates function calls. In contrast, the first subexpression in a special form is *not* evaluated, but rather detected and treated specially by the interpreter. The remaining subexpressions may or may not be evaluated, depending on the behavior of the special form. For instance, `or` will short-circuit when it detects a non-true value, so the above example will not error, since `or` will never reach the divide-by-zero.

- What is the purpose of the `quote` special form?

  The `quote` special form is meant to *postpone* the evaluation of an expression. For instance, if we write `(1 2 3)`, Scheme will typically treat it as a call expression, treating 1 as a procedure (which it is *not*!). Instead, writing
(quote (1 2 3)), or the equivalent shorthand ’(1 2 3), will cause the overall expression to evaluate to the second subexpression of the quote special form, allowing us to obtain (1 2 3) after evaluation, as desired.
2 Scheme Lists

2.1 Draw out a box-and-pointer diagram for the following list:

```
(define nested-lst (list 1 (cons 2 (cons 3 'nil)) '4 5 6) 7))
```

`nested-lst`

Then, write out what Scheme would display for the following expressions:

```
(scm> (cdr nested-lst))
```

```
((2 3) 4 5 6 7))
```

```
(scm> (cdr (car (cdr nested-lst))))
```

```
(3)
```

```
(scm> (cons (car nested-list) (car (cdr (cdr nested-list)))))
```

```
(1 4 5 6)
```
Notice that the builtin \texttt{append} takes in, not a list of lists, but an arbitrary number of lists as arguments, which it then concatenates together. Implement \texttt{better-append}, which behaves in such a manner, allowing the caller to pass in an arbitrary number of arguments. You may use \texttt{concat} from the previous question.

(Hint: look up “variadic functions” in the Scheme spec!)

\begin{verbatim}
(define (better-append . args)
  (concat args))
\end{verbatim}

\begin{verbatim}
scm> (better-append '(1 2 3))
(1 2 3)
scm> (better-append '(1 2 3) '(2 3 4))
(1 2 3 2 3 4)
scm> (better-append '(1 2 3) '(2 3 4) '(3 4 5))
(1 2 3 2 3 4 3 4 5)
\end{verbatim}

Check your understanding

- How can you get the third element of a Scheme list? Draw out a box-and-pointer diagram if you aren’t sure.

  To get the third element of a Scheme list, we need to get the \texttt{car} of the \texttt{cdr} of the \texttt{cdr} of the list - in other words, the third element of \texttt{lst} is \texttt{(car (cdr (cdr lst)))}.

- What is the difference between \texttt{eq?} and \texttt{equal?} in the context of Scheme lists? Construct two lists \texttt{lst1} and \texttt{lst2} such that \texttt{(equal? lst1 lst2)} is \#t but \texttt{(eq? lst1 lst2)} is \#f.

  \texttt{equal?} tests \textit{equality}, and behaves like == in Python - in other words, it returns true if all the corresponding elements of two lists are themselves equal. \texttt{eq?}, in contrast, tests \textit{identity}, and returns true only if its two arguments are in fact the same \textit{object}. Thus, one possibility is simply \texttt{(define lst1 (list 1))} and \texttt{(define lst2 (list 1))}. 